

### Compact Matrix Display

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The following problem arose in the study of the molecular structure of RNA[1]:

One wishes to represent information about the probabilities that various pairs  $(i, j)$  from a sequence of length  $n$  over a finite alphabet occur. It is important to be able to locate accurately from the display which pair is involved and how probable it is on a logarithmic scale. A similar representation problem arises in displaying the strength of connection between pairs of units of neural networks[2].

For large  $n$ , the only compact way to represent the information on a line printer is to encode a different character for each of a finite number of probability levels. The information is then displayed as a matrix. This leads to rather ugly output which is not easily interpretable even if the characters are chosen with the amount of black ink increasing with increasing probability. However, for a first look this is the most efficient way to obtain the information. The problem then is to convert this character output to a high quality image for visual processing.

A succinct way of doing this is by drawing black boxes of varying sizes accurately positioned with lower left corners forming the square matrix of probabilities. T<sub>E</sub>X provides the opportunity to draw such structures by setting sequences of appropriate `\vrules` and to merge such plots with additional text and alphanumeric information. This merging leads as a side effect to the obvious advantage that a complete paper can easily be transferred through the networks by transmitting just a single file. Figure 2 shows that part of the input file which defines the matrix, containing line printer style character data. The related graphic output is shown in Figure 3.

Figure 1 shows the structure of one matrix element, with  $w_o$  ranging from about 5pt to 10pt and  $w_i < w_o$ . The matrix elements can easily be coded as:

```
\hbox to \wo{\vrule height \wi width \wi
             depth 0pt \hfil}      (V1)
```

V1 is working fine for about  $n < 50$  but for larger  $n$  problems arise concerning T<sub>E</sub>X's internal storage ("! T<sub>E</sub>X capacity exceeded ..."), even if the current

page contains nothing other than the matrix and even if  $mem\_max^\dagger$  is set to the maximum of  $2^{16} - 1$ .

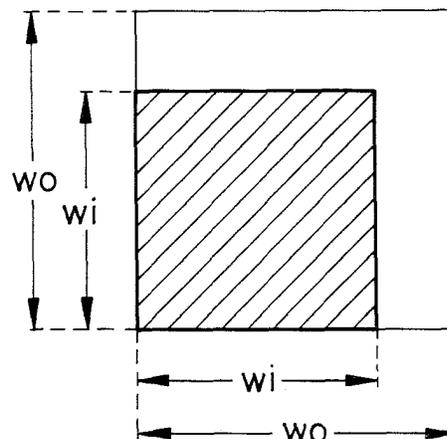


Figure 1

What is the reason for this rapid exhaustion of storage?

Clearly, T<sub>E</sub>X has to hold all the stuff defining the current page in  $mem\_array$  until the `\shipout` operation is done. Taking this into account we have to ask the following question: how much memory does T<sub>E</sub>X use for one matrix element?

Analysing V1 following [3], we find that for each matrix element T<sub>E</sub>X generates the following list of *nodes*:

- 1 *box\_node*
- 1 *rule\_node*
- 1 *glue\_node*

We define  $\sigma$  as the amount of storage T<sub>E</sub>X needs to allocate one rule matrix element. By summing up T<sub>E</sub>X's constants  $box\_node\_size$ ,  $rule\_node\_size$  and  $glue\_spec\_size$  we get:

$$\sigma_1 = 7 + 4 + 4 = 15 \text{ memory\_words}$$

This requirement is high compared with that for a *character token*, which requires only one *memory-word* to fill about the same area on paper.

Fortunately one finds that there is at least one alternative version with  $\sigma < \sigma_1$ :

```
\vrule height \wi width \wi depth 0pt
\rest=\wo
\advance\rest by -\wi
\kern\rest      (V2)
```

<sup>†</sup> terms in italics with enclosed underline character refer to [3]







